

Friday April 5
Lecture 24

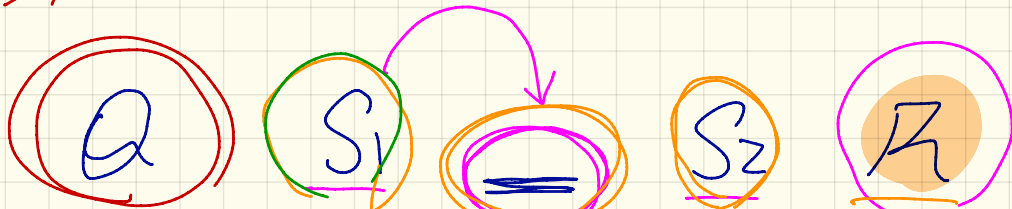
$WP(S_1; S_2, R)$

$= WP(S_1, WP(S_2, R))$ ①

enough for S_2 to establish

R

② upon termination S_1 can be established



$WP(S_2, R)$

intermediate postcondition

Correctness of Program: Sequential Composition

Is $\{ \text{True} \} \text{tmp} := x; x := y; y := \text{tmp} \{ x > y \}$ correct? X

① calculate $wp(\text{tmp} := x; \underbrace{x := y}_{S_1}; \underbrace{y := \text{tmp}; x > y}_{S_2})$

= { wp rule for ; seq. comp. }

$wp(\text{tmp} := x, wp(\underbrace{x := y}_{S_1}, \underbrace{y := \text{tmp}; x > y}_{S_2}))$

= { wp rule for := }

$wp(\text{tmp} := x, wp(x := y, wp(y := \text{tmp}, x > y)))$

$wp(\text{tmp} := x, wp(x := y, x > \text{tmp}))$

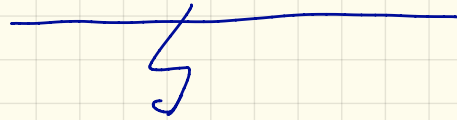
= { wp rule for := }

$wp(\text{tmp} := x, y > \text{tmp}) = \{ wp \text{ rule for :=} \}$
 $y > x$

ex
 $y = 1$
 $x = 2$
 $\text{tmp} \Rightarrow$

$y > x$

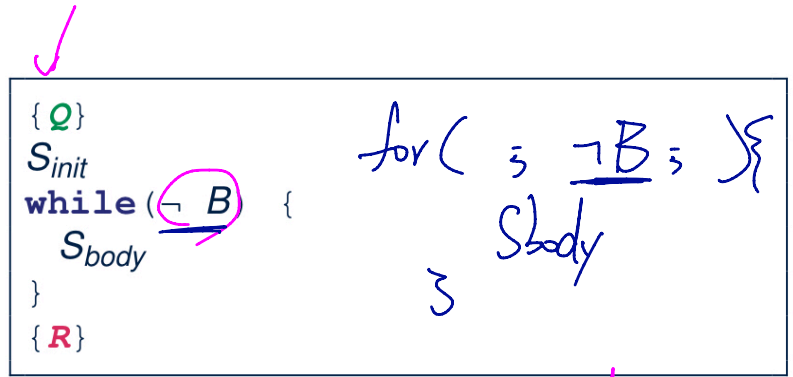
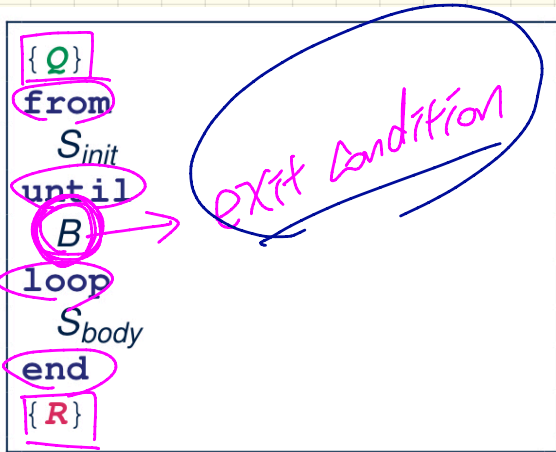
{y > x}



{x > y}

Swap without
introducing an
intermediate variable.

Loops: Eiffel vs. Java

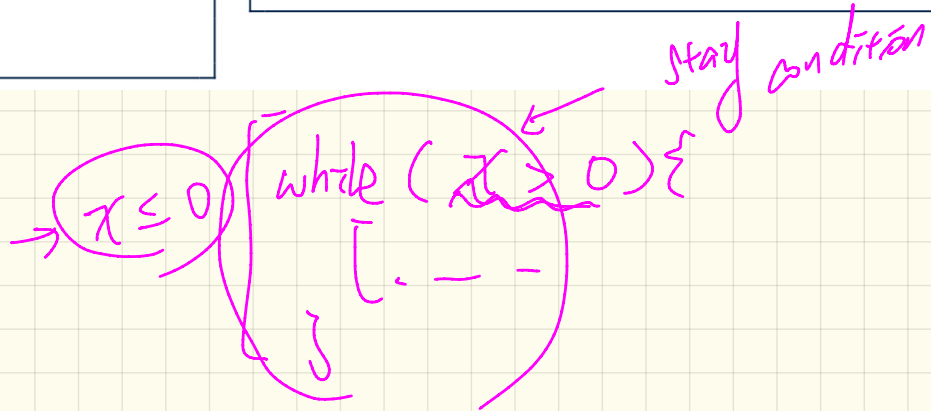


from

until

loop
not ($x > 0$) → $x \leq 0$

end



Contracts of Loops

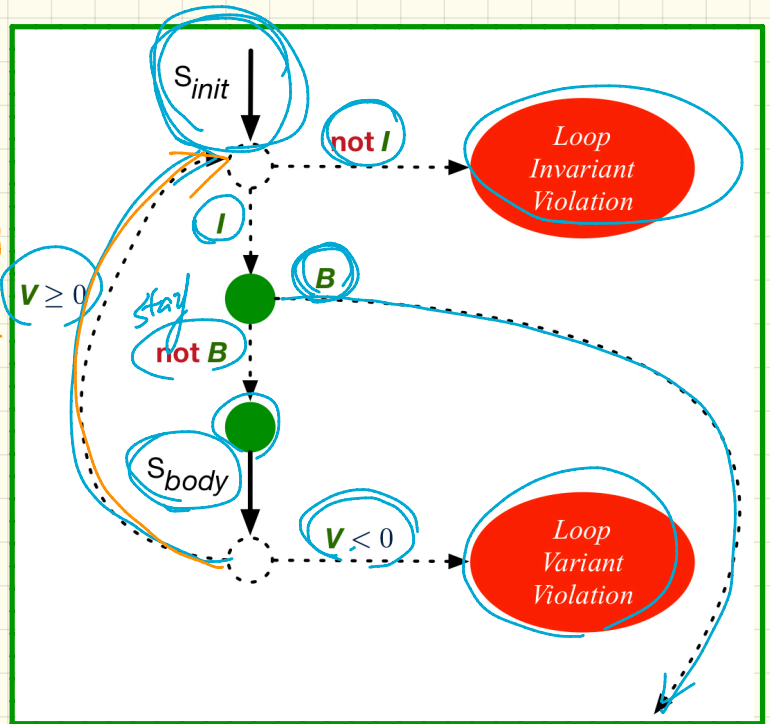
Syntax

```
from
  Sinit
invariant
  invariant_tag: I
until
  B
loop
  Sbody
variant
  variant_tag: V
end
```

established

maintained

Runtime Checks



Contracts of Loops: Example

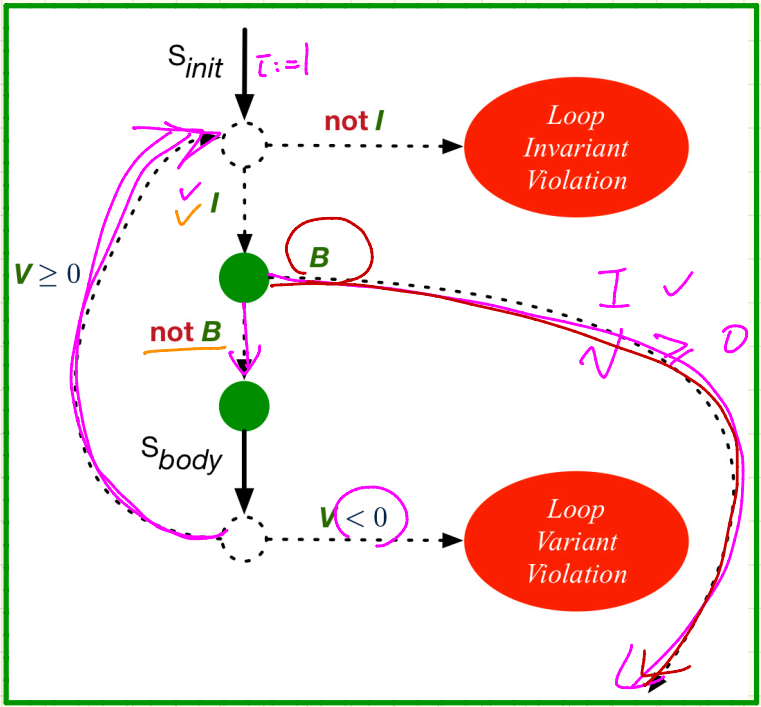
Example

```

test
local
  i: INTEGER
do
  from
    i := 1
  invariant
    1 <= i and i <= 6
  until
    i > 5
  loop
    io.put_string ("iteration " + i.out
    i := i + 1
  variant
    b - 2 = 4
    b - 1
  end
end
end
  
```

Iteration 1 (2)
 Iteration 2
 3
 4
 5 (6)

Runtime Checks



```

test
  local
    i: INTEGER
  do
    from
      i := 1
    invariant
      1 <= i and i <= 5
    until
      i > 5
    loop
      io.put_string ("iteration " + i.out
      i := i + 1
    variant
      6 - i
    end
  end
end

```

← 5
 CI violation
 ∴ after 5th iteration
 E loop ends 6


```

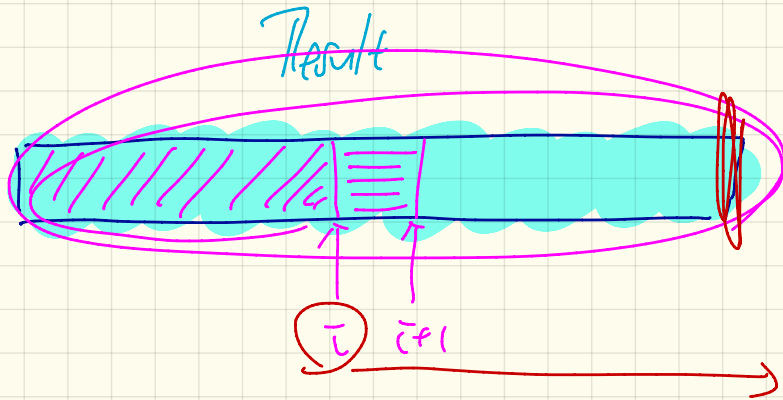
test
  local
    i: INTEGER
  do
    from
      i := 1
    invariant
      1 <= i and i <= 6
    until
      i > 5
    loop
      io.put_string ("iteration " + i.out
        i := i + 1
      variant
        5 - i
      end
    end
  end
end

```

5th iteration

$\bar{2}$ becomes $\underline{6}$

$5 - 6 = (-1)$ LN violation.



Result

Contracts of Loops: Violations

Example

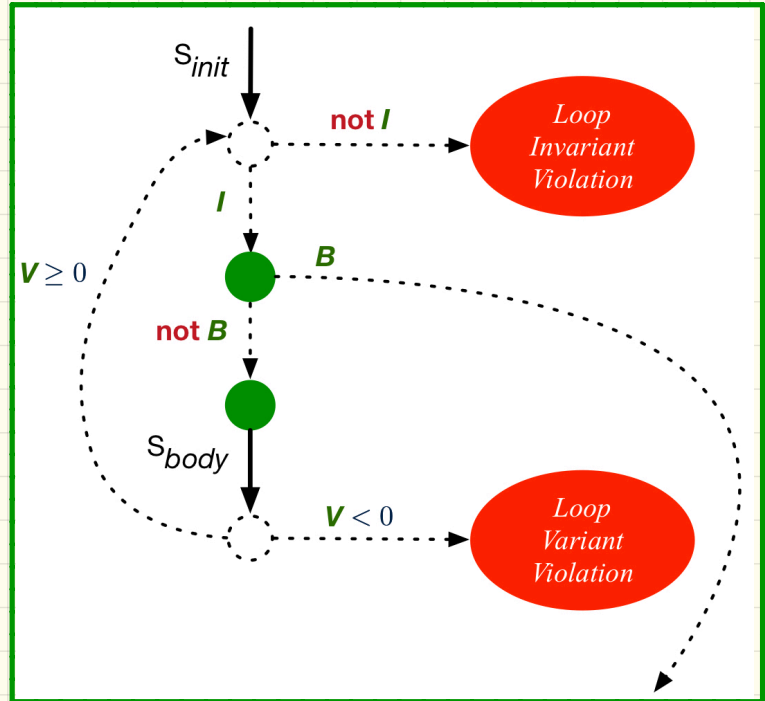
```
test
local
  i: INTEGER
do
  from
    i := 1
  invariant
    1 <= i and i <= 6
  until
    i > 5
  loop
    io.put_string ("iteration " + i.out
    i := i + 1
  variant
    6 - i
  end
end
end
```

Invariant Violation: $1 \leq i \leq 5$

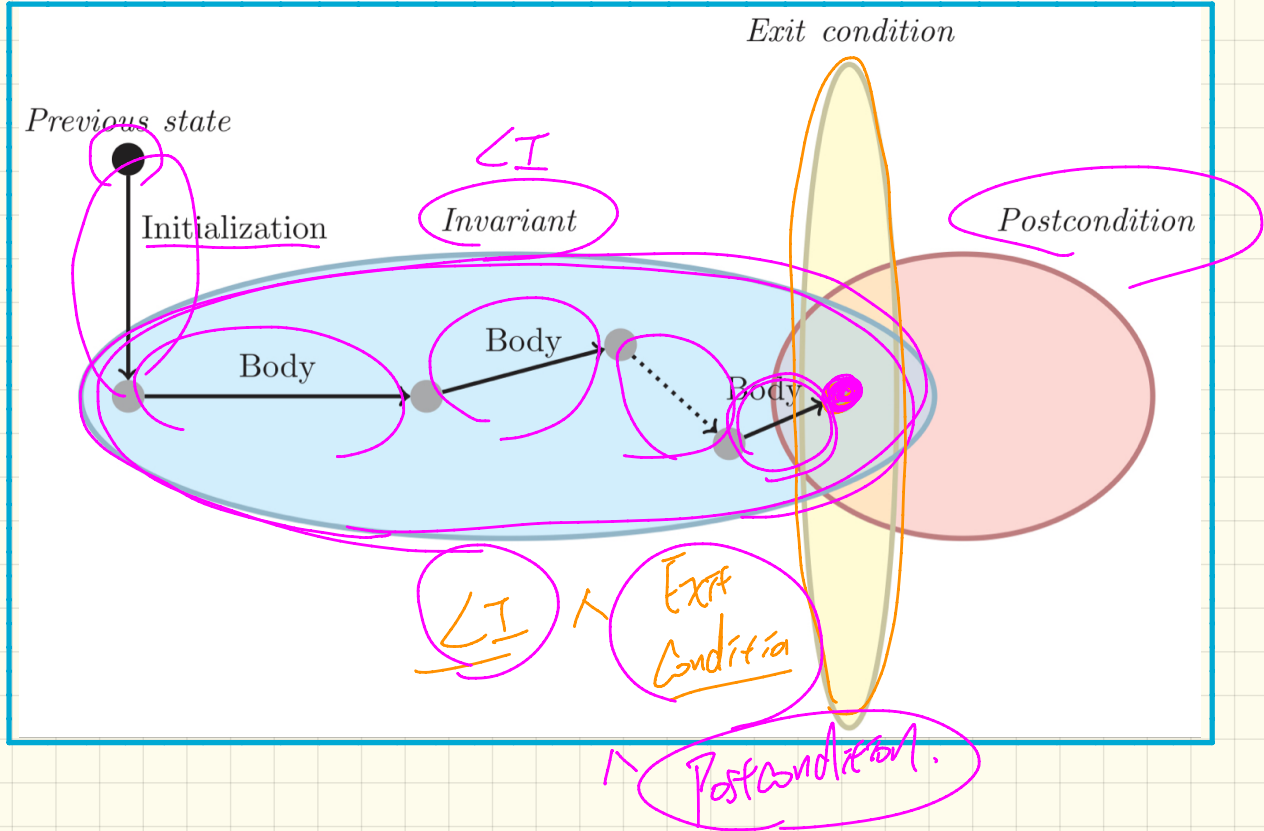
Variant Violation: $5 - i$

Skipping Loop Body: $i > 0$

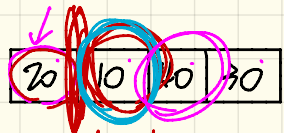
Runtime Checks



Contracts of Loops: Visualization



Finding Max: v1



```

find_max (a: ARRAY [INTEGER]): INTEGER
local i: INTEGER
do
  from
    → i := a.lower; Result := a[i]
  invariant
    loop_invariant: --  $\forall j | a.lower \leq j \leq i \bullet Result \geq a[j]$ 
    → across a.lower |..| i as j all Result >= a [j.item] end
  until
    → i > a.upper
  loop
    → if a [i] > Result then Result := a [i] end
    i := i + 1
  variant
    → loop_variant: a.upper - i + 1
    ensure
      → correct_result: --  $\forall j | a.lower \leq j \leq a.upper \bullet Result \geq a[j]$ 
    → across a.lower |..| a.upper as j all Result >= a [j.item]
  end
end
  
```

$\forall j | a.lower \leq j \leq i \bullet Result \geq a[j]$

i Result
 1 20
 $\forall j | 1 \leq j \leq 1 \bullet 20 \geq a[j]$

is only to be considered in the coming iteration

2 20
 $\forall j | 1 \leq j \leq 2 \bullet 20 \geq a[j]$
 $20 \geq a[1]$
 $20 \geq a[2]$

i
 \rightarrow
 $\forall j | 1 \leq j \leq 3 \bullet 20 \geq a[j]$
 $20 \geq a[3]$

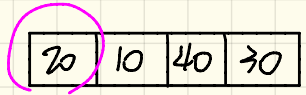
AFTER ITERATION	i	Result	LI	EXIT (i > a.upper)?	LV
Initialization	1	20	✓	×	-
1st	●	●	●	●	●
2nd	●	●	●	●	●

Finding Max: v2

```

find_max (a: ARRAY [INTEGER]): INTEGER
local i: INTEGER
do
  from
    i := a.lower; Result := a[i]
  invariant
  → loop_invariant: --  $\forall j | a.lower \leq j < i \bullet Result \geq a[j]$ 
    across a.lower |..| (i - 1) as j all Result >= a [j.item] end
  until
    i > a.upper
  loop
    if a [i] > Result then Result := a [i] end
    i := i + 1
  variant
    loop_variant: a.upper - i
  end
ensure
  correct_result: --  $\forall j | a.lower \leq j \leq a.upper \bullet Result \geq a[j]$ 
    across a.lower |..| a.upper as j all Result >= a [j.item]
end
end
  
```

i
1 Result
20



F
 $\forall j | 1 \leq j \leq 4$
 $20 \geq a[j]$

AFTER ITERATION	i	Result	LI	EXIT (i > a.upper)?	LV
Initialization	1	20	✓	×	-
1st	2	20	✓	×	2
2nd	3	20	✓	×	1
3rd	4	40	✓	×	0
4th	●	●	●	●	●

$$Vx \mid \underline{I} \cdot P(x) \quad T$$

$$Vx \mid R(x) \cdot P(x)$$

$$\equiv Vx \cdot \frac{R(x)}{I} \Rightarrow P(x)$$

Proof Obligations for Correct Loops

```

{Q}
from
  Sinit
invariant
  I
until
  B
loop
  Sbody
variant
  V
end {R}
  
```

- A loop is **partially correct** if:
 - Given precondition Q, the initialization step S_{init} establishes LI I.

$$\{Q\} S_{init} \{I\}$$
 - At the end of S_{body} , if not yet to exit, LI I is maintained.

$$\{I \wedge \neg B\} S_{body} \{I\}$$
 - If ready to exit and LI I maintained, postcondition R is established.

$$I \wedge B \Rightarrow R$$
- A loop **terminates** if:
 - Given LI I, and not yet to exit, S_{body} maintains LV V as non-negative.

$$\{I \wedge \neg B\} S_{body} \{V \geq 0\}$$
 - Given LI I, and not yet to exit, S_{body} decrements LV V.

$$\{I \wedge \neg B\} S_{body} \{V < V_0\}$$

↓ $B \wedge I$

$V \geq 0$